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“Preamble-Free Frame Synchronization for Short-Packet Low-Latency Communications”

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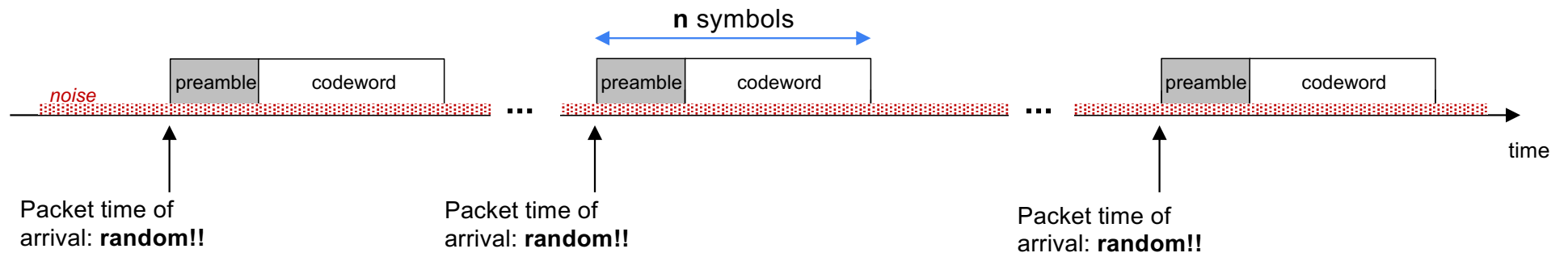
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Problem Statement

Asynchronous packet-based digital link:



Low Latency → **SHORT packets**
($n < 100$)

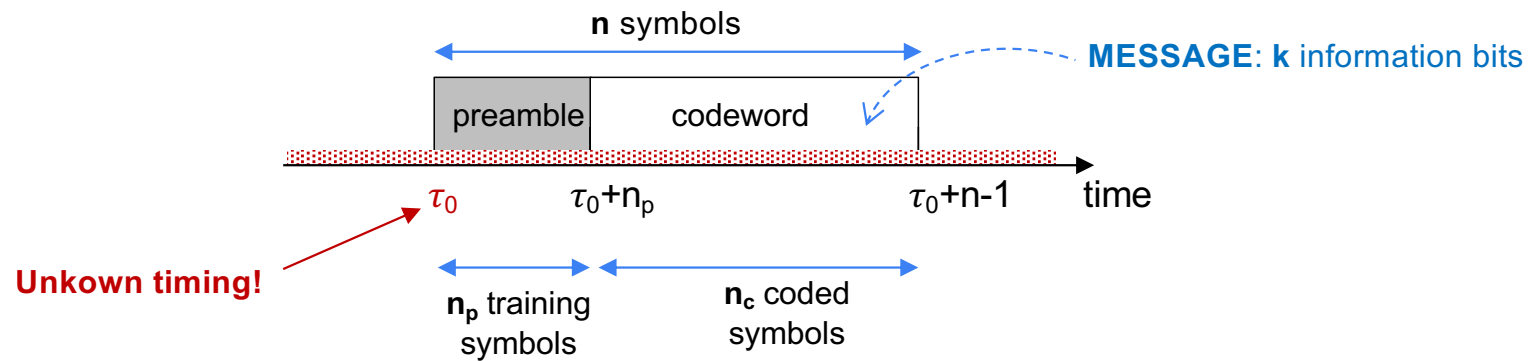
SHORT preambles !!
Synchronization ??

Code-aided (CA) synchro

SHORT codewords !!
Reliability ??

Short-blocklength codes
(Polar, TB-CC, ...)

Problem Statement



OBJECTIVE:

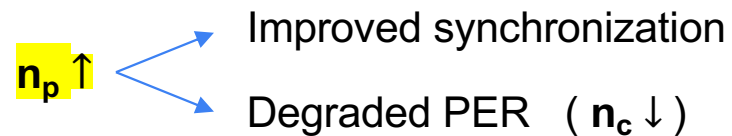
Message decoding with minimum PER $\rightarrow$ **requires estimation of τ_0**

n : **fixed** (latency constraint)

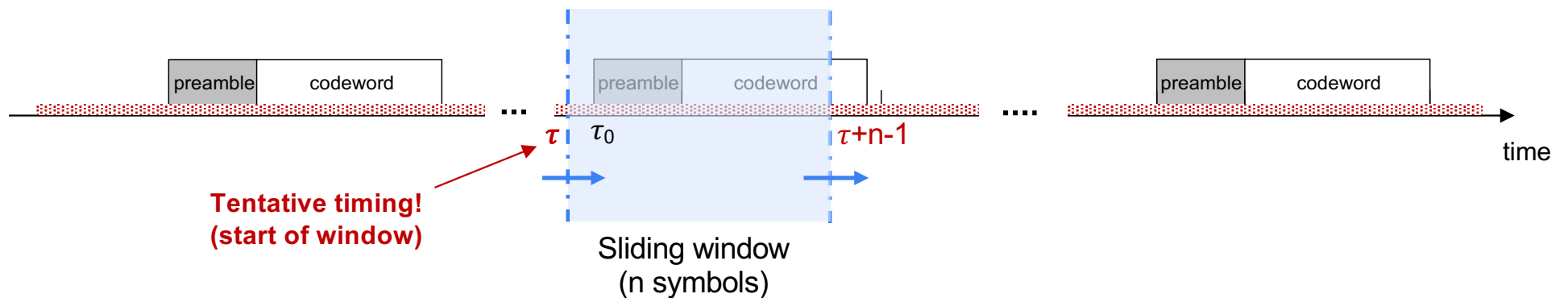
k : **fixed** (bit rate constraint)

n_p : to be optimized

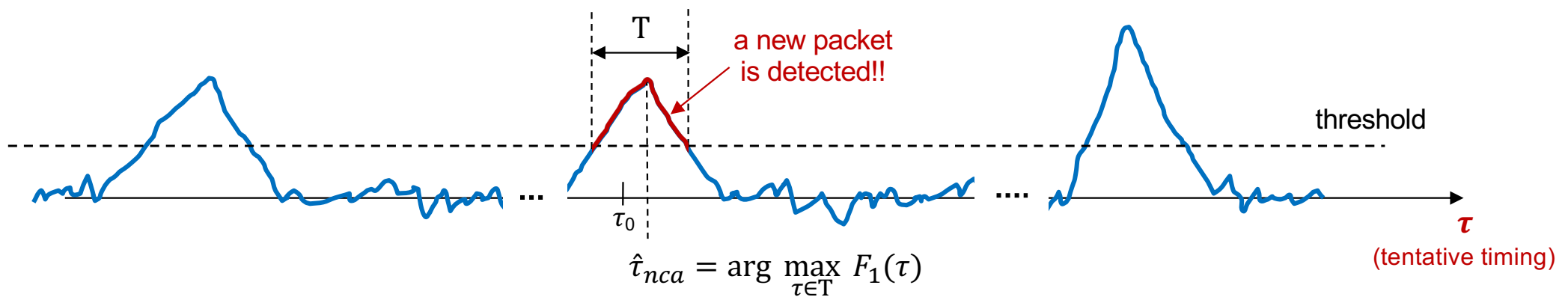
$$n_c = n - n_p$$



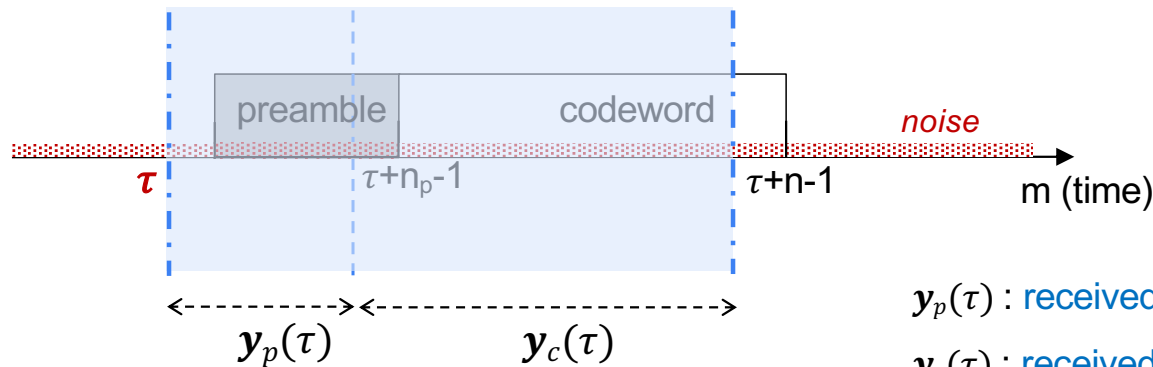
Synchronization (step 1)



Non-Code-Aided (NCA) synchronization metric $\rightarrow F_1(\tau)$



Synchronization (step 1)



$\mathbf{y}_p(\tau)$: received preamble samples assuming τ (length $\mathbf{n}_p$)

$\mathbf{y}_c(\tau)$: received codeword samples assuming τ (length $\mathbf{n}_c$)

(at the matched filter output)

Non-Code-Aided (NCA) synchronization metric :

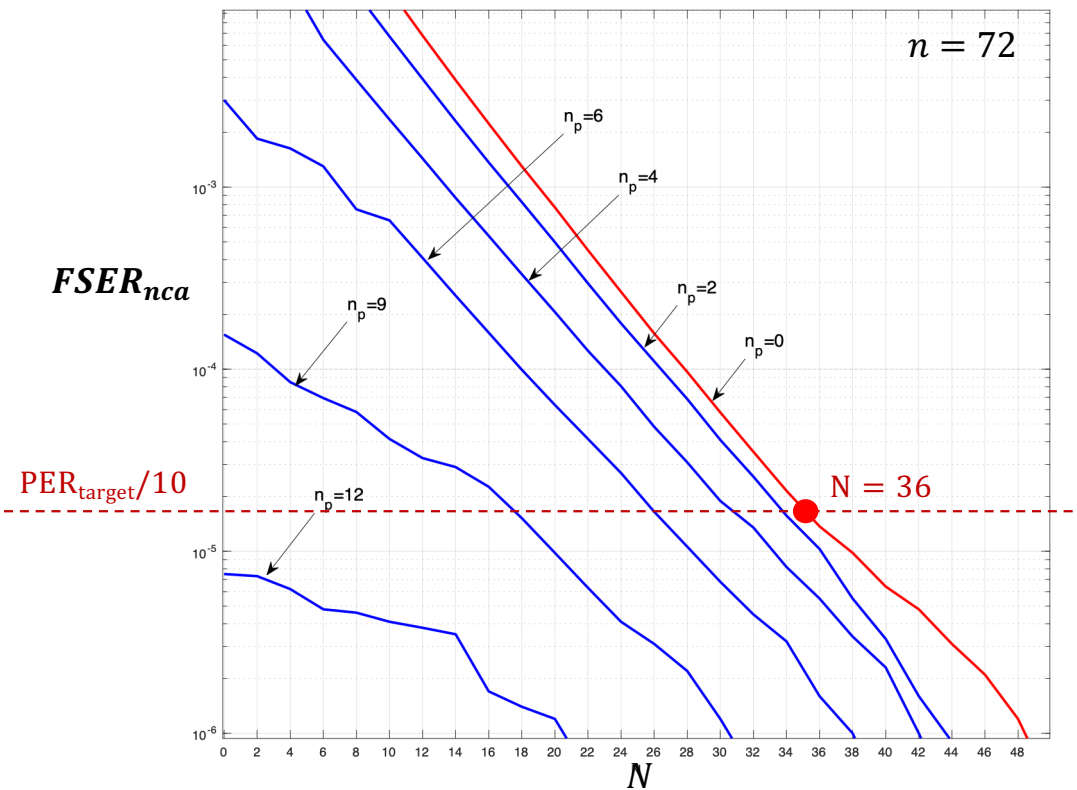
$$F_1(\tau) \doteq \frac{2}{\sigma_w^2} \Re \{ \mathbf{y}_p^H(\tau) \mathbf{x}_p \} + \sum_{m=0}^{n_c-1} \ln \cosh \left(\frac{2\Re\{[\mathbf{y}_c(\tau)]_m\}}{\sigma_w^2} \right) + \sum_{m=0}^{n_c-1} \ln \cosh \left(\frac{2\Im\{[\mathbf{y}_c(\tau)]_m\}}{\sigma_w^2} \right)$$

Preamble
(Data-Aided)
QPSK-modulated Codeword
(Non-Data-Aided)

Classical Massey's metric [24] adapted for burst transmission (Maximum Likelihood criterion). Good approximation: $\ln \cosh(x) \approx \|x\|_1$

Synchronization (step 1)

Non-Code-Aided (NCA) synchronization performance analysis:

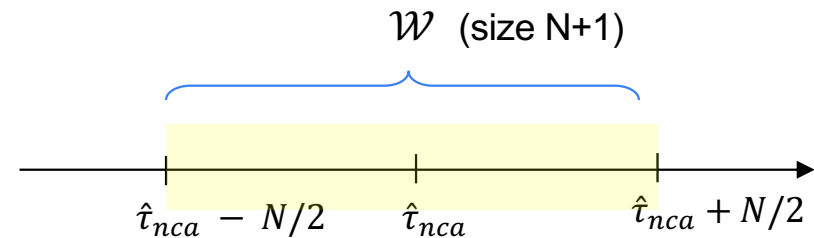


Frame Synchronization
Error Rate

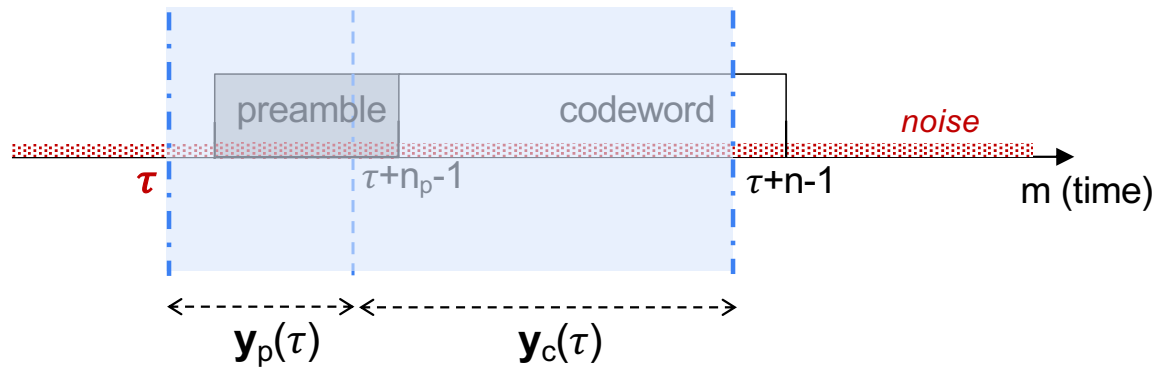
$$FSER_{nca} = \text{Prob} \left(|\hat{\tau}_{nca} - \tau_0| > \frac{N}{2} \right)$$

$$= \text{Prob} (\tau_0 \notin \mathcal{W}) < PER_{target}/10$$

Negligible!!



Synchronization (step 2)



Code-Aided (NCA) synchronization metric :

$$F_2(\tau) \doteq \mathfrak{R}\{\mathbf{y}_p^H(\tau)\mathbf{x}_p\} + \mathfrak{R}\{\mathbf{y}_c^H(\tau)\mathbf{f}_{\text{enc}}[\hat{\mathbf{b}}(\tau)]\}$$

Preamble
(Data-Aided)
Codeword
(Code-Aided)

Joint Maximum Likelihood (ML)
synchronization and decoding.

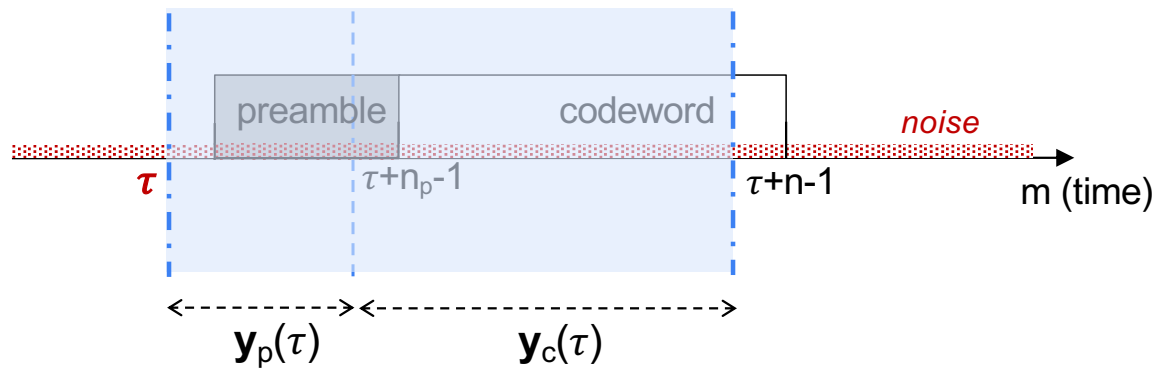
for $\tau \in \mathcal{W}$

Reduced search
window!!



$$\hat{t}_{ca} = \arg \max_{\tau \in \mathcal{W}} F_2(\tau)$$

Synchronization (step 2)

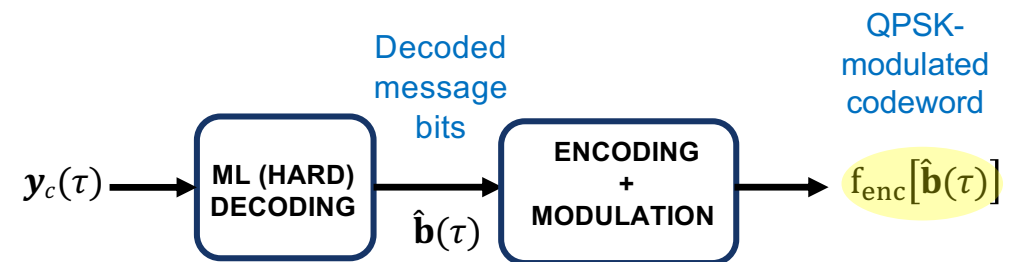


Code-Aided (NCA) synchronization metric :

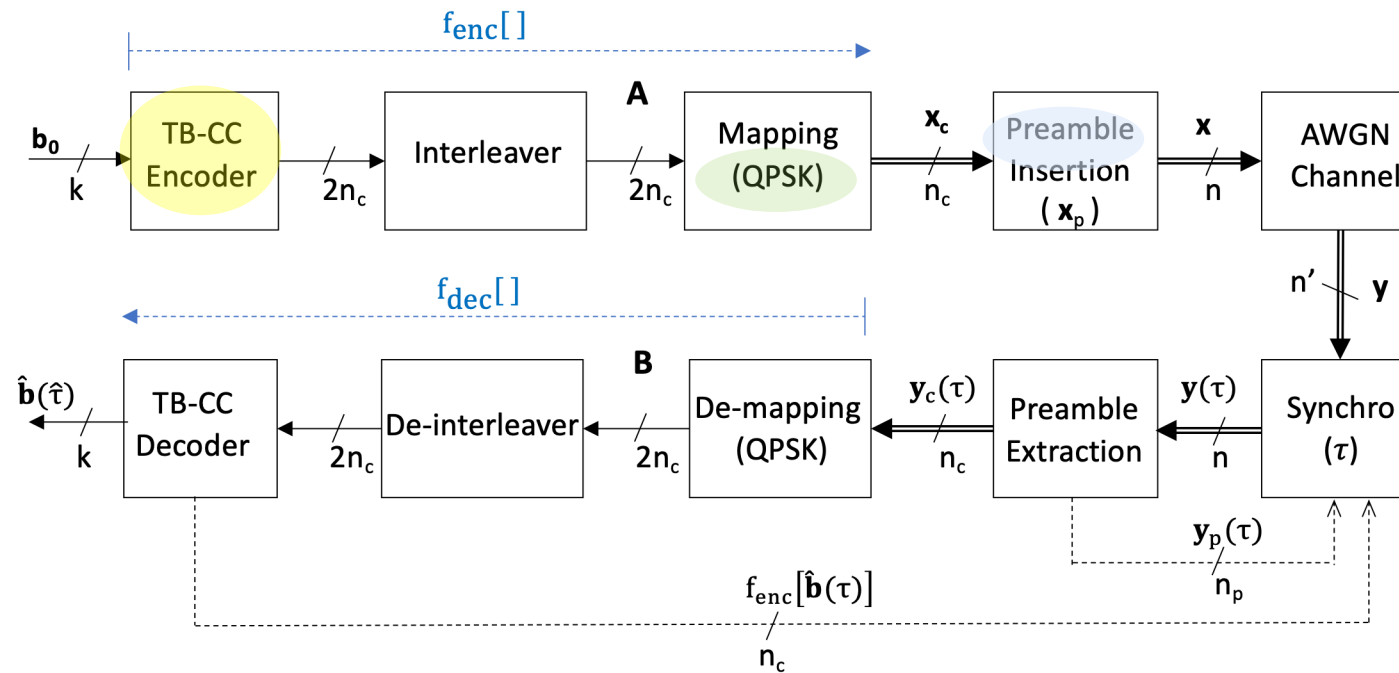
$$F_2(\tau) \doteq \mathfrak{R}\{\mathbf{y}_p^H(\tau)\mathbf{x}_p\} + \mathfrak{R}\{\mathbf{y}_c^H(\tau)\mathbf{f}_{\text{enc}}[\hat{\mathbf{b}}(\tau)]\}$$

Preamble
(Data-Aided)
Codeword
(Code-Aided)

Joint Maximum Likelihood (ML)
synchronization and decoding.



Overall architecture

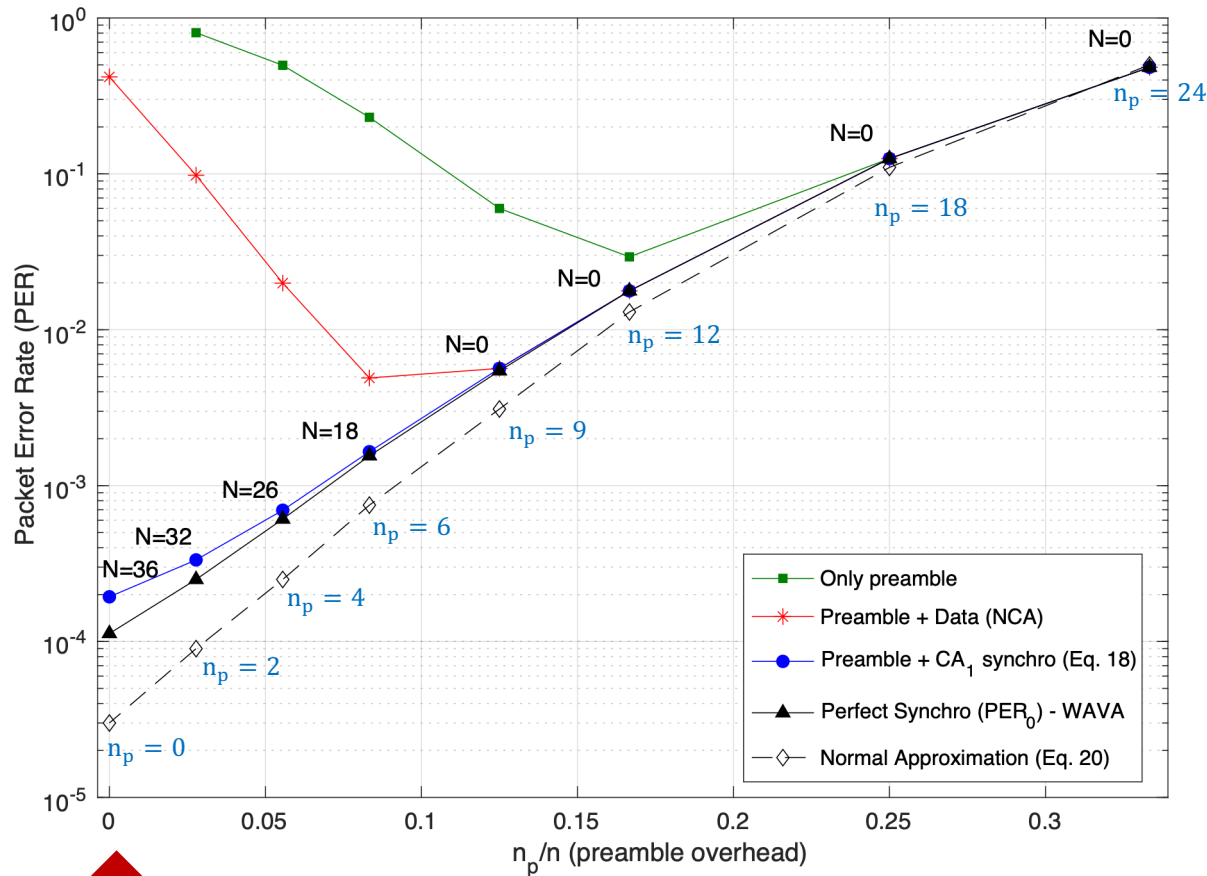


Code: TB-CC (large-memory Tail-Biting Convolutional Codes)

Modulation: QPSK

Preamble: complex-valued **Zadoff-Chu** sequence

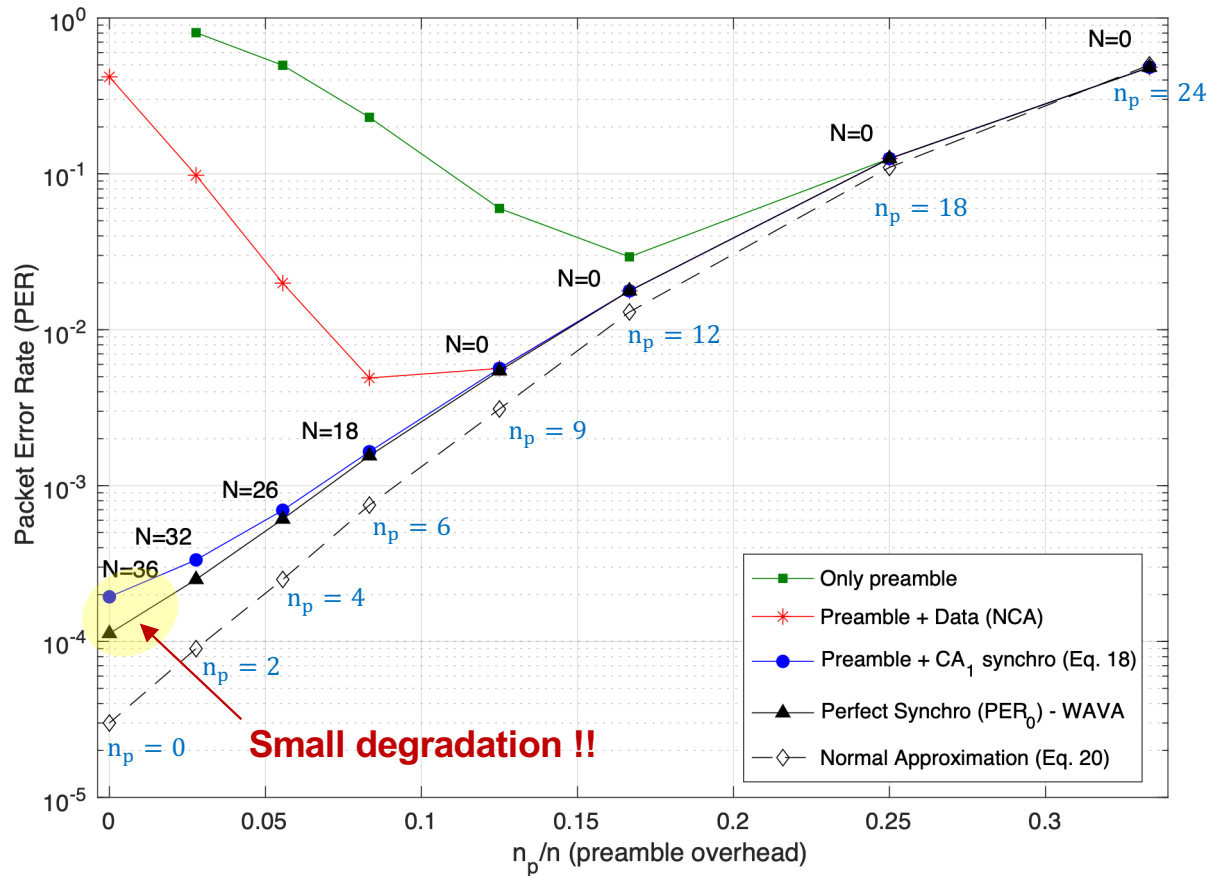
Simulations



Preamble-less

- Packet length: $n = 72$ (fixed!)
- Message length: $k = 72$ (fixed!)
- **TB-CC**
 $\frac{1}{2}$ -rate mother code: $5537_{\text{oct}} \ 6131_{\text{oct}}$
 Memory = 11 bits (2048 states)
 Variable code rate (puncturing): $R_c = \frac{k}{2(n-n_p)}$
 Wrap-Around Viterbi decoder (WAVA)
 (iterations ≤ 4)
- AWGN channel ($E_b/N_0 = 3\text{dB}$)
- Correct carrier phase/frequency synchronization
- Correct symbol (fine timing) synchronization.

Simulations (II)



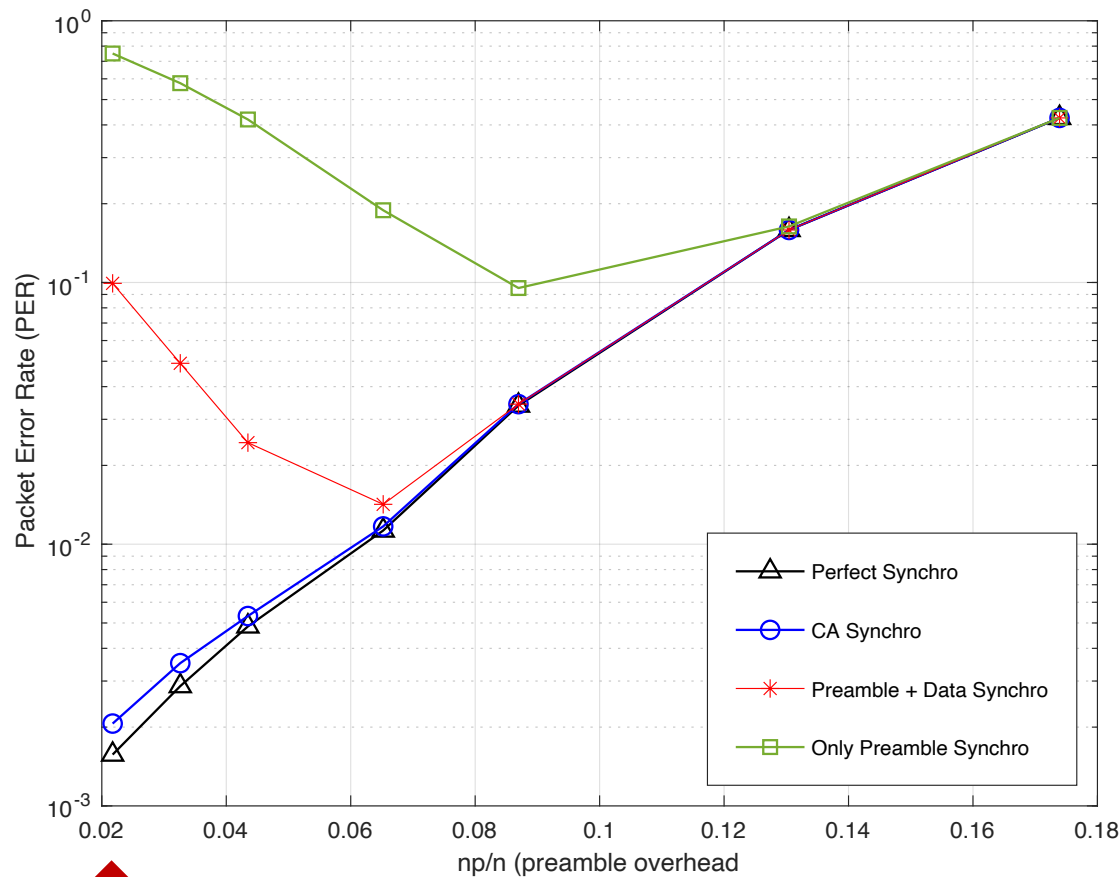
Code Memory	FSER	PER_0	PER
4	1.8×10^{-5}	2.6×10^{-2}	2.6×10^{-2}
8	6.8×10^{-5}	1.0×10^{-3}	1.0×10^{-3}
11	8.0×10^{-5}	1.2×10^{-4}	1.8×10^{-4}
14	8.4×10^{-5}	5.0×10^{-5}	1.2×10^{-4}

Synchronization does not improve (actually degrades) if we increase the TB-CC memory.



For short packets and large code memory, synchronization degrades the overall PER.

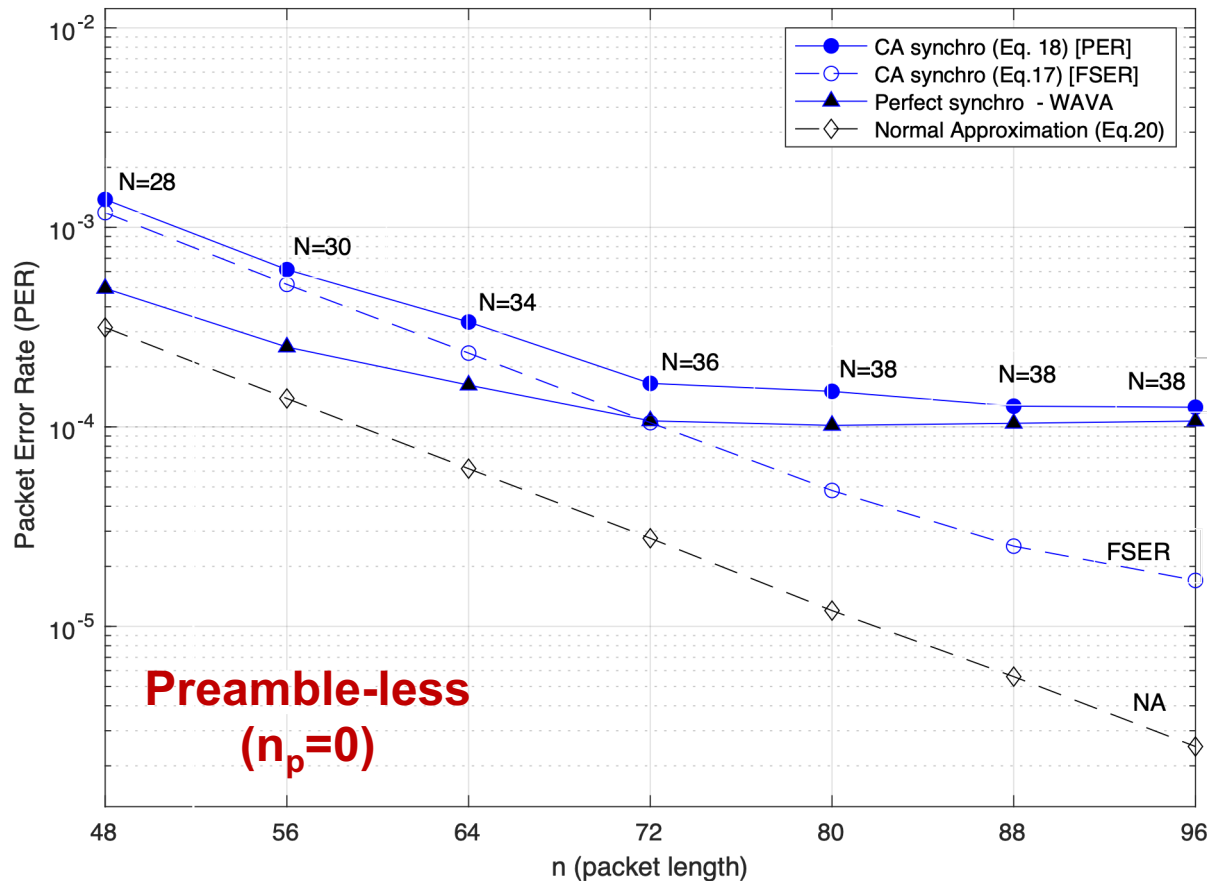
Simulations (III)



↑
Preamble-less

- Packet length: **$n = 48$ (fixed!)**
- Message length: **$k = 48$ (fixed!)**
- **TB-CC**
 $\frac{1}{2}$ -rate mother code: 5537_{oct} 6131_{oct}
 Memory = 11 bits (2048 states)
 Variable code rate (puncturing): $R_c = \frac{k}{2(n-n_p)}$
 Wrap-Around Viterbi decoder (WAVA)
 (iterations ≤ 4)
- AWGN channel ($E_b/N_0 = 3\text{dB}$)
- Correct carrier phase/frequency synchronization
- Correct symbol (fine timing) synchronization.

Simulations (IV)



- Packet length: $n = 48-96$ (variable!)
- Message length: $k = n$
- **TB-CC**
 $\frac{1}{2}$ -rate mother code: 5537_{oct} 6131_{oct}
 Memory = 11 bits
 No puncturing: $R_c = 1/2$ (fixed!!)
 Wrap-Around Viterbi decoder (WAVA)
 (iterations ≤ 4)
- AWGN channel ($E_b/N_0 = 3\text{dB}$)
- Correct carrier phase/frequency synchronization
- Correct symbol (fine timing) synchronization.

Conclusions and further work

- TB-CC are suited for short blocklength communications ($n < 72$). Gap to Normal Approximation decreases for short packets ($n \leq 72$) and large memory ($m \geq 11$)
- Synchronization of TB-CC encoded packets can degrade the overall PER in that case (i.e, $n \leq 72$ and $m \geq 11$).
- **Preamble-less is optimal** in terms of PER for fixed packet and message lengths (n and k fixed) even for very short packets (e.g., $48 \leq n \leq 72$).
- Preamble-less communications requires **$N+1$ decodings** per packet. For $E_b/N_0=3\text{dB}$ and $n=k=72$, TB-CC, we found that **$N \geq 36$** .
- **Complexity reduction** requires:
 - To reduce the number of decodings per packet ($N+1$). Parallelisation to not increase latency!
 - To reduce complexity of decodings